The neglected art of Fixed Point arithmetic

Jetro Lauha Seminar Presentation

Assembly 2006, 3rd - 6th August 2006 (Revised: September 13, 2006)

2019-08, small note: Many of the statements in this presentation do not hold true for "today's hardware". (floating point support is now common in mobile and even IoT CPUs).

Contents

- Motivation
- Introduction
- Typically needed functions
- Caveats and tricks
- Tips for making a fixed point library

- Man sent himself to moon, and space probes even beyond that. Do you think the hardware used to accomplish those feats had fancy FPU to do all the calculations?
- They used RCA 1802.
 - Processing power equals roughly 6502 or 6510, used in Apple II and Commodore 64.

- But it's a lot of work.
 - 30% of the Apollo software development effort was spent on scaling. [KrL64]
- So they eventually switched to floating point when hardware got better.

- So why am I talking about this?
 - Well, at least it's COOL, in retro-way:
 This is how demo & game coders did their 3D stuff
 15 years ago and made some pretty cool stuff even with the minuscule CPU power.
- But does that matter anymore except if you are going to take part in the old school demo competition with some retro stuff?

- There's still plenty of platforms where using only fixed point (integer) calculations is still very relevant.
 - Mobile devices (Typical: ARM CPU, no FPU)
 - Almost all mobile phones (J2ME or native code)
 - Handheld consoles (Gameboy, Nintendo DS)
 - DSP Programming
 - There's both fixed & floating point DSPs

- ...continued...
 - OpenGL ES is the standard for embedded 3D.
 - Profiles for both fixed point and floating point, but often only Common-Lite profile is provided (no floating point).
 - Fixed point is often still a bit faster on desktop than floating point.
 - Stable calculations across platforms
 - Floating point calculations are prone to slight differences based on compiler, CPU and other dependencies.

Introduction

- Basics
- Notation
- Range and precision
- Conversion
- Basic operations: + * /

Introduction: Basics

- What are the fixed point numbers in "layman's" terms?
 - Scale all real numbers by a constant factor, such as 65536, round to nearest integer and and store the numbers as integers.
 - This allows you to represent an evenly distributed subset of real numbers roughly from -32768 to 32767 (with 32-bit signed integers and factor of 65536).

Introduction: Basics

• More exactly, you are dividing your range of values to two parts - the integer part and fractional part.



Introduction: Notation

- Notations:
 - *M*.*N*, e.g. 16.16
 - QN (Q factor), e.g. Q16
- *M* is number of integer bits and *N* is number of fractional bits.

Introduction: Range and precision

• Range: defined by the integer (upper) part.

- 16.16 (signed): range is [-32768, 32767]

 Precision: smallest difference between two successive numbers is 1/2^N.

- 16.16: 1/65536 (~0.000015258789)

8-bit example:

4.4 Range: [-8, 7] (if signed) Precision: 1/16 (0.0625)

Introduction: Conversion

- Conversion from real to fixed point number
 - Multiply by 2^{N} and round to nearest integer
 - (int) (**R** * (1<<**N**) + (**R**>=0 ? 0.5 : -0.5))
- Conversion from fixed point to real number
 - Cast to real and divide by $2^{\scriptscriptstyle N}$
 - (float)**F** / (1<<**N**)
- Conversion from/to integers (lossless)
 - Shift N bits up or down (scaling by 2^{N})

•
$$F = I << N$$
, $I = F >> N$

Introduction: Basic operations

- Addition (+) and subtraction (-)
 - Same as adding and subtracting integers
- Multiplication (a * b)
 - Multiply as integers and divide result by 2^{N} .
 - ((a * b) >> N)
 - That overflows very easily, as both a and b are fixed point numbers!
 - If both a and b are 2.0 (131072) as 16.16 fixed point (a * b) == 17179869184 - 32 bits isn't enough!

Introduction: Basic operations

- For multiplication, the intermediate result from (a * b) is in 2M:2N (Q2N) format
 - Store intermediate value in double sized integer format. That is, for 32-bit 16.16 fixed point numbers, you need a 64-bit integer to store the 32.32 (Q32) intermediate result.

<u>INT64</u>

MSVC: __int64 GCC: signed long long Java: long

Introduction: Basic operations

- Division (a / b)
 - Multiply a by 2^{N} and divide by b (as integers).

• ((a << N) / b)

- Again, intermediate result is prone to overflowing, so the correct way for 16.16 is:
 - (((INT64)a << N) / b)
- See references for more detailed introductory texts to fixed points. [VVB04, Str04, WikF]

Typically Needed Functions

- Sine and cosine: sin(x), cos(x)
- Arcus tangent: atan2(y, x)
- Square root: sqrt(x)
- Try CORDIC

Typically Needed Functions: Sine and cosine

- Typical approach is to use a look-up table.
 - Requires memory proportional to desired accuracy
 - Requires some storage space to load table from or time for pre-calculating table on startup
 - Can interpolate between sampled values to gain some more accuracy
- Note that it's enough to calculate $\pi/4$ entries to table, rest of the samples can be mirrored and transformed from those.

Typically Needed Functions: Sine and cosine

• It's possible to find or construct less accurate approximations for functions if you need smaller code, memory usage or more speed.

- DSP coders have some quite nice tricks. [Ben06]

• See also [Str04] for code example of how to calculate sin, cos and tan algorithmically using only a small arctan table.

Typically Needed Functions: Square root

- Several fairly good iterative algorithms exist, so I don't recommend using a look-up table.
- Can be as simple as trying out to multiply integers by themselves until you find out the closest one
 - Or binary search version of the above
- Ken Turkowski's implementation is probably the most often used one. [Tur94]
 - For your convenience, code on the next slide.

Typically Needed Functions: Square root

```
/* The definitions below yield 2 integer bits, 30 fractional bits */
#define FRACBITS 30 /* Must be even! */
#define ITERS (15 + (FRACBITS >> 1))
typedef long TFract;
TFract
FFracSgrt (TFract x)
{
   register unsigned long root, remHi, remLo, testDiv, count;
   root = 0;  /* Clear root */
remHi = 0;  /* Clear high part of partial remainder */
   do {
       remHi = (remHi << 2) | (remLo >> 30); remLo <<= 2; /* get 2 bits of arg */
       root <<= 1; /* Get ready for the next bit in the root */
       testDiv = (root << 1) + 1;  /* Test radical */</pre>
       if (remHi >= testDiv) {
           remHi -= testDiv;
           root += 1;
   } while (count-- != 0);
   return(root);
                                                                        [Tur94]
```

Typically Needed Functions: Arcus tangent

- You can try some look-up table tricks, again.
- If fast and rough approximation is enough, implementation can be very simple. [Cap91]
- For accurate results, try using CORDIC (covered next).
- For my favorite approximation (for the time being), check Jim Shima's DSP Trick: Fixed-Point Atan2 With Self Normalization. [Shi99]

Typically Needed Functions: Try CORDIC

- "COordinate Rotation DIgital Computer", an algorithm to calculate hyperbolic and trigonometric functions, from 1959. [WikC]
 - Only small look-up tables, bitshifts and additions.
- Use it run-time or to pre-calculate look-up tables. (sin, cos, atan, ...)
- Accurate results
- Not the fastest solution

Caveats And Tricks

- Back to range and precision
- Watch out for division by zero
- Exact results
- Dealing with problems

Caveats And Tricks: Back to range and precision

- When storing result of a*b to normal sized fixed point (integer) value
 - Possible range & precision for the original values is much more limited than the normal to prevent overflow & underflow.
 - For storing a*a:
 - abs(a)<=~181 -- 181*181 = 32761, barely fits in signed 16.16 fixed point number.
 - abs(a)>=~0.004 -- 0.004*0.004 = 0.000016, truncated down to 1/65536.

Caveats And Tricks: Back to range and precision

- Similarly, make sure that a/b will stay in range
 - When |b| > 1.0
 - Check ranges so that result doesn't end up being 0.
 - When |b| < 1.0
 - b>1/(2^{M-1}/a)
 - If max value for a is 32, b must be at least 0.000991821 (65/65536) so that a/b fits in 16.16 fixed point number: 32/0.000991821=~32263.
 - If b would be one less (64/65536), then a/b will be 32768, not fitting in [-32768, 32767] 16.16 fixed point value range.

Caveats And Tricks: Watch out for division by zero

- Floating points have "Infinity Arithmetic"
 - Even result of division by zero is defined, so you simply get Inf as a result
 - Easier to go unnoticed by mistake
- Fixed point (integer) division by zero leads to interrupt or an exception is thrown
 - Typically programs just crash at this

Caveats And Tricks: Exact results

- Possible in some cases: modify division involving formulas to keep numerator and denumerator separate, and try to find out final (exact) result by examining those, without doing the division. See [Eri05] for example.
- Generally speaking, it's rare and hard to take advantage of this.

Caveats And Tricks: Dealing with problems

- When troubled by overflows, underflows or accuracy problems
 - Try keeping the intermediate result(s) in the bigger (64 bit) format and work out the final result directly from there.
 - Use asserts and do other verification checks rigorously, especially in debug builds.
 - Compare to results of same calculations done in floating points.

Tips For Making A Fixed Point Library

- There's built-in support... if you code in Ada.
- C/C++ alternatives:
 - Code it all in-line, using normal integers
 - Use helper macros (conversions, operations)
 - Create a real number class with overloaded operators
 - Allows to switch easily between floats and fixed points

Tips For Making A Fixed Point Library

- Create debug version of the real number class
 - Perform both fixed point and floating point calculations in parallel
 - Detect overflow & underflow conditions
 - Detect drifting
 - Error/warning asserts and checks can be made run-time togglable
- If you work on J2ME, it's best to inline all calculations yourself for performance.

Other Tidbits

- Nobody noticed that I changed the underlying physics engine from floating point to fixed point in latest version of *Pogo Sticker*.
- You can do fixed point (integer) abs() without branches. [And05, War02]
 - For 32-bit ints:
 - result = (v ^ (v >> 31)) (v >> 31)
 - Ridiculously that's patented. But that's not the only way, check the references.

Other Tidbits

- 32-bit signed 0x8000000 (highest bit) is special
 - int x; if (x < 0) x = -x; Doesn't work as expected if x==0x8000000! X will still be 0x8000000 (-2147483648).
 - For the above example, solution is to cast result to unsigned int as you know it will not be negative.

References

KrL64	http://www.hq.nasa.gov/office/pao/History/computers/Ch4-2.html – Kreide, H., Lambert, D.W., <i>Computation: Aerospace Computers in Aircraft, Missiles and Spacecraft</i> , Space/Aeronaut., 42, 78 (1964); see also N.H. Herman and U.S. Lingon, <i>Mariner 4 Timing and Sequencing</i> , Astronaut. Aeronaut., 43 (October 1965).
VVB04	Van Verth, J. M., Bishop, L. M., <i>Essential Mathematics for Games & Interactive Applications – A Programmer's Guide</i> , Morgan Kaufmann, 2004.
Str04	Street, M., A Fixed Point Math Primer, OpenGL [®] ES Game Development, Course Technology PTR, 2004.
WikF	http://en.wikipedia.org/wiki/Fixed-point_arithmetic – Fixed-point arithmetic article in Wikipedia.
Tur94	Turkowski, K., <i>Fixed Point Square Root</i> , Apple Technical Report No. 96, 1994. Also appears in <i>Graphics Gems V</i> , Paeth, A. W. (editor), Academic Press, 1995. See http://www.graphicsgems.org/.
Cap91	Capelli, R., <i>Fast Approximation to the Arctangent</i> , Graphics Gems II, Academic Press, 1991. See http://www.graphicsgems.org/.
Shi99	http://www.dspguru.com/comp.dsp/tricks/alg/fxdatan2.htm – Shima, J., <i>DSP Trick: Fixed-Point Atan2 With Self</i> <i>Normalization</i> , post in comp.dsp newsgroup, Apr 23, 1999.
Ben06	http://www.audiomulch.com/~rossb/code/sinusoids/ – Bencina, R., Fun with Sinusoids, 2006.
WikC	http://en.wikipedia.org/wiki/CORDIC – CORDIC article in Wikipedia, see especially the referenced "CORDIC Bibliography Site" and the C implementation by Peter Knoppers (http://people.csail.mit.edu/hqm/imode/fplib/cordic_code.html).
Eri05	Ericson, C., <i>Numerical Robustness for Geometric Calculations (aka EPSILON is NOT 0.00001!)</i> , GDC Proceedings, 2005. Also available from http://realtimecollisiondetection.net/pubs/.
And05	http://graphics.stanford.edu/~seander/bithacks.html – Anderson, S. E., Bit Twiddling Hacks.
War02	Warren, H. S., <i>Hacker's Delight</i> , Addison-Wesley, 2002.

Thank You!

URL for these slides:

https://iki.fi/jetro/2006/08/07/neglected-art-of-fixed-point-arithmetic/

Fill out this form if you're interested in more information about Fixed Point Math:

https://docs.google.com/forms/d/e/1FAIpQLScZ56aEt7oJED-kDFFlaUHJZ6FLy3AZ520P9gHYMv8OAtIsVg/viewform

• Short URL: http://j.mp/morefixedpoint